# Welcome to CS 3000: Algorithms \& Data! 

Section 1<br>Instructor Tim LaRock (he/him/his)<br>larock.t@northeastern.edu<br>bit.ly/cs3000syllabus

## Zoom Notes

I will be recording our Zoom lectures.

Keep both your video and audio muted at all times unless you are speaking.

- Multiple video streams increases the bandwidth required for a smooth video.
- As I understand it:
- If you are muted, you are not part of the recording.
- If you unmute your video or audio, you will be recorded.

If you have a question, use the chat box to either (a) write your question directly or (b) indicate you would like to ask a question out loud.

- I prefer the chat to the "raise hand" feature because it is persistent.

The Zoom chat is always archived. I will probably delete it very soon after recording.

## Today

## Brief instructor introduction

Some presentation of the what/why of Algorithms

Course logistics + questions

Some content

## Me

## Tim LaRock (he/him/his)

## Just call me Tim!

I grew up in the Adirondack Mountains


Researcher at the Network Science Institute
Usually: Understanding how things move through networks, e.g. how a ship moves through a network of ports.

Lately: Analyzing mobility data to understand the impact of mobility restrictions on the spread of COVID-19.

Now: Your instructor!

## This Course

We are going to learn about algorithms, which are sets of instructions for how to manipulate data

Erickson definition: "An algorithm is an explicit, precise, unambiguous, mechanically-executable sequence of elementary instructions, usually intended to accomplish a specific purpose."

Specifically, we will cover things like...

- Transforming problems from informal descriptions to formal mathematical descriptions
- Formulating strategies for solving formal problems efficiently
- Understanding, designing, and choosing appropriate data structures for our solutions
- Proving the correctness of a solution mathematically
- Determining the complexity in terms of (i) running time and (ii) memory requirements for a proposed solutions
- Categorize problems and solutions based on classes of complexity
- ...and much more!

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3. How we can go from the input to a solution?
4. Can we guarantee that a solution is correct?
5. Can we guarantee a solution is found in a reasonable amount of time?
6. And more...

In this course, we learn mathematical techniques that allow us to effectively communicate answers to these questions.

## Reason 2: Efficient algorithms are important in practice!

Scalability or efficiency of an algorithm can be the difference between a computation running in 5 minutes or never finishing before the heat death of the universe.

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"Recipe" is a classic example of an algorithm


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Algorithm for constructing 1 PB\&J sandwich


Input: 2 slices of bread, jar of PB, jar of jelly, spreading tool

Algorithm:

1. Use the tool to spread PB on one slice of bread
2. Use the tool to spread jelly on the slice of bread without peanut butter
3. Put the two slices of bread together so that the PB and J are facing each other.
4. Cut in half if desired.

Output: PB\&J

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Algorithm for constructing 1 PB\&J sandwich


Input: 2 slices of bread, jar of PB, jar of jelly, spreading tool, desire, direction

Algorithm:

1. Use the tool to spread PB on one slice of bread
2. Use the tool to spread jelly on the slice of bread without peanut butter
3. Put the two slices of bread
together so that the PB and J are facing each other.
4. Cut in half if desired.

CutInHalf(desire, direction)
Output: PB\&J

## Reason 2: Efficient algorithms are important in practice!

Scalability or efficiency of an algorithm can be the difference between a computation running in 5 minutes or never finishing before the heat death of the universe.

Assume it takes 2 minutes to make a sandwich.

Algorithm for constructing 1 PB\&J sandwich


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Assume it takes 2 minutes to make a sandwich.

Algorithm for constructing 1 PB\&J sandwich

What if I want $\mathrm{N} \gg 1$ PB\&J sandwiches?

Input: 2 slices of bread, jar of PB, jar of jelly, spreading tool

Algorithm:

1. Use the tool to spread PB on one slice of bread
2. Use the tool to spread jelly on the slice of bread without peanut butter
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Output: PB\&J

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Scalability or efficiency of an algorithm can be the difference between a computation running in 5 minutes or never finishing before the heat death of the universe.

Assume it takes 2 minutes to make a sandwich.

Algorithm for constructing N PB\&J sandwiches.

Runtime: N*2 minutes

## What if I want N >> 1

## PB\&J sandwiches?

## REPEAT N TIMES:

Input: 2 slices of bread, jar of PB, jar of jelly, spreading tool

## Algorithm:

1. Use the tool to spread PB
on one slice of bread
2. Use the tool to spread jelly
on the slice of bread without peanut butter
3. Put the two slices of bread together so that the PB and J are facing each other.
4. Cut in half if desired.

Output: PB\&J

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Scalability or efficiency of an algorithm can be the difference between a computation running in 5 minutes or never finishing before the heat death of the universe.

Runtime: $\mathrm{N}^{*} 2$ minutes

Probably fine if I want less than $\mathrm{N}=10$ sandwiches.

If I want $\mathrm{N}=1000$, I will quickly run out of resources and time!

## What if I want $\mathrm{N} \gg 1$

## PB\&J sandwiches?

## REPEAT N TIMES:

Input: 2 slices of bread, jar of PB, jar of jelly, spreading tool

## Algorithm:

1. Use the tool to spread PB
on one slice of bread
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3. Put the two slices of bread together so that the PB and J are facing each other.
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## Reason 2: Efficient algorithms are important in practice!

Scalability or efficiency of an algorithm can be the difference between a computation running in 5 minutes or never finishing before the heat death of the universe.

Sometimes we don't even know if a scalable solution to a problem could possibly exist - the techniques you learn here will give you the tools to answer that question!

# Reason 3: Algorithms/complexity theory is an interesting field of mathematics 

Theoretical advances have serious practical implications ( $\mathrm{P}=\mathrm{NP}$ )


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## Bonus Reason

Studying algorithms often feels like solving a puzzle!


## Logistics



## Logistics - Course Structure

Lectures (like this one) Monday - Thursday, 1:30-3:10PM
Homework - Approximately weekly (45\% of grade)
Exams - 2 Midterms (15\% each) and a final exam (25\%)

## Resources:

Course website: bit.ly/cs3000syllabus
Canvas: Contact me ASAP (larock.t@northeastern.edu) if you do not have access!

## Logistics - Lecture details

Lectures (like this one) Monday - Thursday, 1:30-3:10PM
Recorded live and uploaded to Canvas
Attendance is encouraged if possible, but not required

Regardless of live attendance, it is expected that you have watched the lecture at some point! Anything discussed in lecture is fair game for homework/exams!

## Logistics - Homework details

Assigned approximately weekly, with variation depending on timing of exams
Not meant to take you hours upon hours to complete - if you are stuck, ask for help (more to come on various ways to do so)

Collaboration is okay!
Write solutions in your own words and include all collaborators names on everyone's submissions
Copying is not okay!!
We reserve the right to ask you to explain any answer you submitted!
Okay to look up resources online for help, but..
ALWAYS evaluate your sources carefully!
A textbook page is preferable to Wikipedia, Wikipedia is much more reliable than a stack exchange answer with 0 votes, etc. Use your judgement!

NEVER copy solutions if you find them.
If you find an exact answer and can't "unsee it", do not copy it! Just send me an email.

## Logistics - Exam details

Exams - 2 Midterms and a final
All "take home" format, meaning you will have a set time period to work on them outside of class

Similar to homework assignments, except absolutely NO collaboration is allowed and use of the internet is limited to textbooks ONLY

## PLEASE DO NOT CHEAT!

Obviously we are on the honor system, and my default attitude is to trust you! But if you are caught cheating there will be severe penalties, including escalation to the College and/or University level.

## Logistics - Instructor Office Hours

I will hold open office hours at the following times:

> 4-5 PM on Tuesdays
> 8:30-9:30 AM on Wednesdays
> 1:30-2:30 PM on Fridays

The specific structure of these hours is not decided and depends somewhat on level of demand.

You can always reach out via email to schedule a 1-1 or small group conversation with me.

## Logistics - TAs and Office Hours

## We have 8 Teaching Assistants for the course - they are a resource!

| Name | email | Office Hour 1 | Office Hour 2 |
| :--- | :--- | :--- | :--- |
| Saurabha | jirgi.s@husky.neu.edu | Wednesday, 10AM-11AM | Thursday, 11AM-12PM |
| Ronn | jacob.r@husky.neu.edu | Wednesday, 12PM-1PM | Thursday, 12PM-1PM |
| Himanshu | budhia.h@husky.neu.edu | Tuesday, 4PM-5PM | Monday, 12PM-1PM |
| Dania | abuhijleh.d@husky.neu.edu | Monday, 9AM-10AM | Wednesday, 9AM-10AM |
| Drew | bodmer.d@husky.neu.edu | Thursday, 10AM-11AM | Monday, 3PM-4PM |
| Angela | gross.an@husky.neu.edu | Wednesday, 12PM-1PM | Friday,Friday 2PM-3PM |
| Luke | boyer.l@husky.neu.edu | Monday, 7PM-8PM | Tuesday, 8PM-9PM |
| Kevin | hui.k@husky.neu.edu | Wednesday, 6PM-7PM | Thursday, 6PM-7PM |

## Canvas

Northeastern's replacement for Blackboard

New to me, new to $1 / 3$ of you (according to entry form)

Plan to use it for a couple of things:
Assignment submission
Grades
Online discussions

Everything else will be at: bit.ly/cs3000syllabus

## Logistics - Online Discussion

Canvas has online discussion boards, I encourage you to post there when you have questions you aren't ready to bring to me/the TAs yet!

Obviously it is not okay for anyone to post solutions on the discussion board, but clarifying and helping guide classmates is okay.

## Textbooks

I will assign some reading from two freely available books:

1. Algorithms by Jeff Erickson
2. Introduction to Algorithms by Cormen, Leiserson and Rivest (CLR)

See the syllabus and Canvas for links, or just search for the titles.

Algorithm Design by Tardos and Kleinberg is no longer required

- If you got a copy, it is a great resource that I encourage you to use!


## Answers to entry form questions

50 people filled it out - thank you!

## Answers to entry form questions

- Will this course require a 0 -credit recitation like in the fall, and will we also need to take this?
- To my knowledge, there is no 0-credit recitation for this course.
- Approximately how many homework assignments will we have in this class?
- Between 5-7, with lowest grade dropped.
- How many hours should I take per day to compete the homework?
- Homework assignments are not meant to take over your life. I expect between 1-5 hours over the course of a week (e.g. $\sim 1 \mathrm{hr}$ per day) to be enough.
- When are homework assignments typically due?
- Still working out the full schedule, but most will likely be due either Fridays or Mondays.


## Answers to entry form questions

- When will we know the dates/times for midterms/finals?
- I will try to finalize the schedule and let you know ASAP.
- Is there any coding in this class?
- No, sorry. I may ask a coding question (in your fav language) as extra credit, but this class is not intended to teach you about programming.
- Can I get some useful websites for learning LaTeX?
- Yes. I will also try to spend a few minutes in one of these early lectures talking about the very basics.
- A couple of people expressed frustration with assumed knowledge across courses.
- I will try to explain concepts as I introduce them, but there are prerequisites for the course so some familiarity with the topics covered in those is assumed! Feel free to ask in the chat if you aren't sure what I am talking about.

More questions before we move to content?

## Content

## Refresher: What is an asymptote?

"...an asymptote (/'æsimptovt/) of a curve is a line such that the distance between the curve and the line approaches zero as one or both of the $x$ or $y$ coordinates tends to infinity." - Asymptote on Wikipedia

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Hor does your
Asymptotes and Runtime algorithen Scale
with N?
linew in $n$ what do asymptotes have to do with algorithms? $O\left(n^{\phi}\right)$
 $N$ is the "size" of log (1) the inapt to $a_{n}$ algocithon

## Sorting

Sorting is extremely important to computer users and scientists!

## Sorting

## Constant time is written

Sorting is extremely important to computer users and scientists!

A simple example: Finding the median of a set of numbers

```
Input: L, an array of N numbers
Output: The median of L
Procedure:
    1. Sort L
    2. If N is odd, return the number at L[\lceil\frac{N}{2}\rceil]
    3. If }\textrm{N}\mathrm{ is even, return the mean of the
        numbers at L[\lceil\frac{N}{2}\rceil] and L[\lceil\frac{N}{2}\rceil+1]
```


## Bubble Sort

Idea: Items "bubble up" to the top as they are sorted pairwise


## Bubble Sort



Idea: Items "bubble up" to the top as they are sorted pairwise

```
Input: L, an array of }\textrm{N}\mathrm{ numbers
Output: L sorted in ascending order
Procedure:
    Let swapped = True
    while swapped = True:
        swapped = False
        for i from 1 to N-1:
            if L[i] > L[i+1]:
                        Swap L[i] and L[i+1]
            swapped = True
```

Bubble Sort Example
Scrapped is TAal
53622

$$
\begin{array}{ll}
35622 & \begin{array}{l}
3 \\
35226
\end{array} \\
3256 & \text { scrapped }==\text { Tree to The } \\
322565 & \text { swapped }=\text { True } \\
22356
\end{array}
$$

Bubble Sort Analysis
loop-invarinanf


Running Time
Best cause: $O(n)$
Worst Cun:O( $\left.n^{2}\right)$

## Next Time

A better approach to sorting

Divide and Conquer Algorithms

More asymptotic analysis

Suggested Reading:

- Erickson book: Introduction thru Chapter 1.1
- CLR book: Introduction thru Chapter 2

Homework 1: To be released tomorrow, due next Monday

