Lecture 7: Dynamic Programming

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bit.ly/cs3000syllabus

Business

Homework 2 is out, due Tuesday May 19 11:59PM Boston time on Canvas

We are working on grading homework 1, will share solutions once the grades are complete

It is totally fine to look ahead in the slides, but please give others a minute to try to answer questions I ask before writing what you saw later

Today

Dynamic Programming

- Fibonacci Numbers
- Text Segmentation Revisited

Fibonacci Numbers

$$f_n \begin{cases} 0 & if\ n = 0 \\ 1 & if\ n = 1 \\ f_{n-1} + f_{n-2} & otherwise \end{cases}$$

Fibonacci Numbers: Recursion

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Fibonacci Numbers: Recurrence Relation

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$$T(n) = 2T(n-1) + 1 \leq 2 \cdot 2^n$$

$$\leq O(2^{n+1})$$

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$$T(n) = 2f_{n+1} - 1 \rightarrow 2T(n+1) \leq O(2^{n+1})$$

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$$T(n) = 2f_{n+1} - 1 \rightarrow 2T(n+1) \leq O(2^{n+1})$$

Exponential in n is *very* slow for such a simple function!

Memoization

```
Fib(n):
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Memo(r)ization

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- What if instead we save each value we compute so that we can access it in constant time?

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  If n = 0:
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```

```
MemFib(n):
  If n = 0:
    return 0
  ElseIf n = 1:
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  Else:
    If F[n] is undefined:
      F[n] = MemFib(n − 1) + MemFib(n − 2)
    return F[n]
```

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- What if instead we save each value we compute so that we can access it in constant time?
- Keep a global table $F[i]$ that stores results and use stored results where possible

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- What if instead we save each value we compute so that we can access it in constant time?
- Keep a global table $F[i]$ that stores results and use stored results where possible
- How is the table filled? And what implication does this have for the runtime?

Memo(r)ization

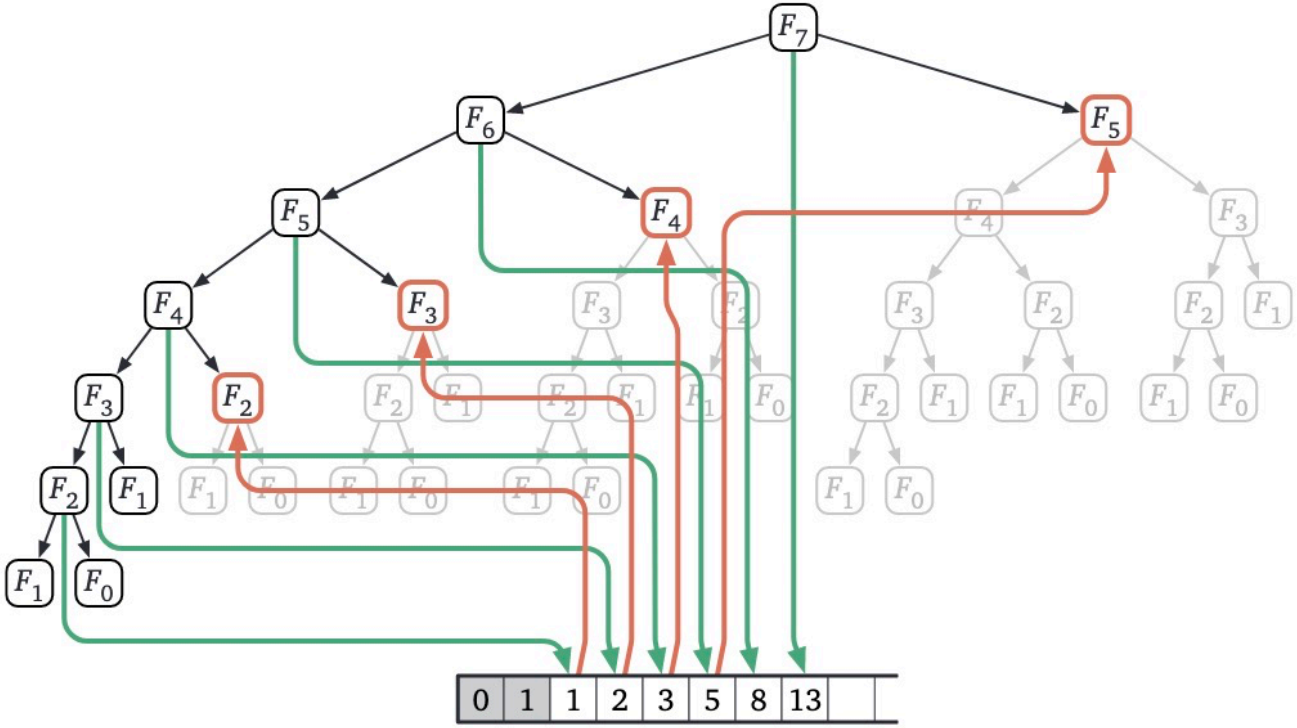

Figure 3.2. The recursion tree for F_7 trimmed by memoization. Downward green arrows indicate writing into the memoization array; upward red arrows indicate reading from the memoization array.

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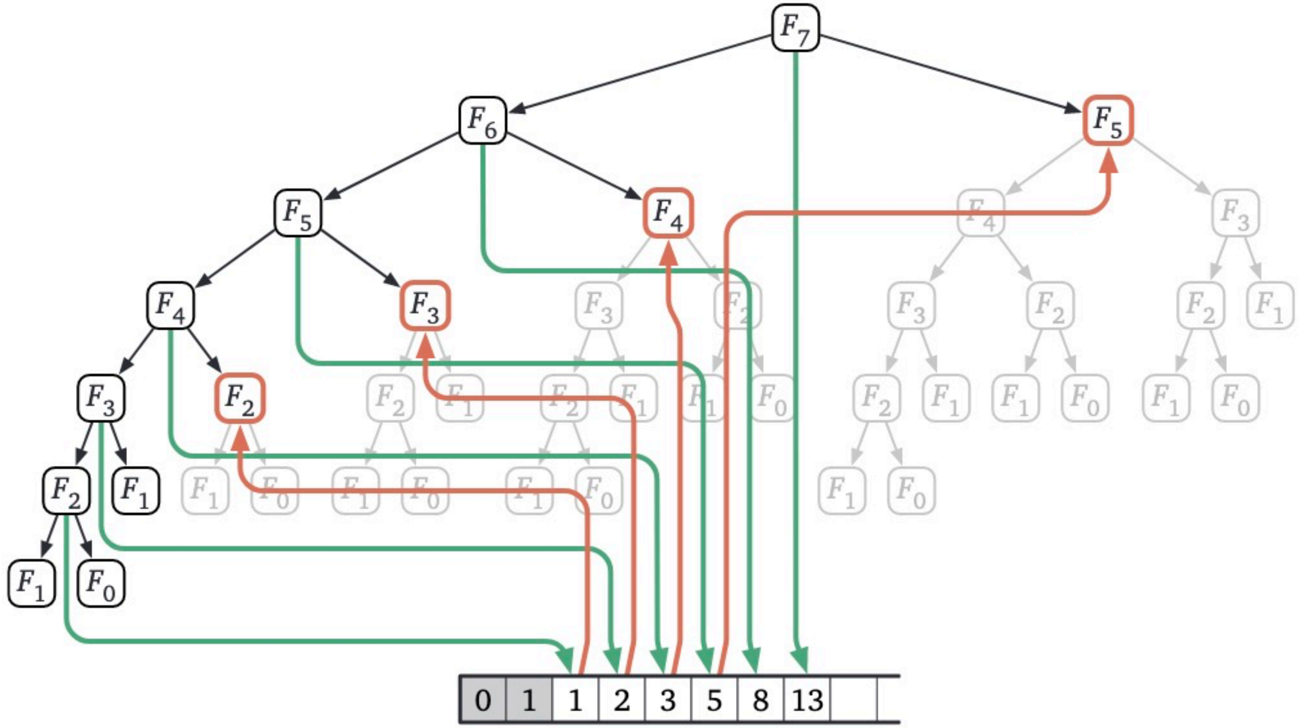
Memo(r)ization

How many additions?



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    return $F[n]$

There has to be a better way!

Enter: Dynamic programming

Dynamic Programming

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- The execution order and runtime of $MemFib(n)$ implies a simpler way to compute Fibonacci numbers

Dynamic Programming

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```
IterFib(n):
  F[0] ← 0
  F[1] ← 1
  for i from 2..n
      F[i] ← F[i − 1] + F[i − 2]
  return F[n]
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- What if we just like….filled F explicitly?

Dynamic Programming

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- The execution order and runtime of $MemFib(n)$ implies a simpler way to compute Fibonacci numbers
- What if we just like….filled F explicitly?
- Now the execution is clearly $O(n)$!
- Note: We could save memory here. How?

Dynamic Programming

- Formalized by Richard Bellman at RAND in the ‘50s
  - Bellman apparently named it “dynamic programming” to obscure his research from his bosses.
  - Programming does not refer to computers, but scheduling: for example designing the “program” of a performance or event, or filling a TV schedule

- General pattern: Recursion without repetition
  - Store solutions of intermediate problems to be reused later!
  - Finding a correct recurrence that can be memoized is vital
    - If your recurrence is wrong or can’t be memoized, you will go in circles!

Dynamic Programming Process

There are 3 main steps to developing dynamic programming solutions:

1. Find the right recurrence
   - Formalize the problem carefully
   - Find a recursive solution (could be pseudocode or just a relation)

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   - What are the subproblems that need solving?
   - What data structure can I use to access them correctly and quickly?
   - Which subproblems depend on each other?
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3. Prove it!

Text Segmentation Revisited

$T(n) = 2T(n-1) + c \ \leq O(2^n)$

```
Splittable(A[1..n], i):
  If i > n:
    return True
  Else:
    j ← i
    for j to n:
      If IsWord(i,j):
        If Splittable(A[1..n], j + 1):
          return True

  return False
```

Problem: Given an array $A[1..n]$ representing a sequence of n characters without spaces, determine whether the array can be subdivided into a sequence of *words.*

Assume we are given a function $IsWord(i,j)$. This function assumes A is a global variable and returns True if the subarray $A[i..j]$ is a word in the language of the sequence.

- This allows us to avoid passing subarrays as arguments to functions.

Text Segmentation Revisited

Where are we wasting computation?

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$O(n)$

For all indices $1 \le \text{i} \le j \le n$, how many times can we call $IsWord(i, j)$?

$O(n^2)$

We are spending exponential time computing polynomial amounts of stuff!

Dynamic Programming Approach

$FastSplittable(A[1..n])$:
   $SplitTable[n+1] \leftarrow True$

   for i from n to 1:
      $SplitTable[i] \leftarrow False$
      for j from i to n:
         If $IsWord(i,j)$ AND $SplitTable[j+1]$:
            $SplitTable[i] \leftarrow True$

   return $SplitTable[1]$

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$n = 18$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k

$i = 18$

$j = 18$

SplitTable

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k	
F	F	F	F	F	F	F	F	F	F	F	F	F	F	F	F	F	F	T

$j + 1 = 19$

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$j + 1 = 18$

Dynamic Programming Approach

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$SplitTable$

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   $SplitTable[n+1] \leftarrow True$

   for i from n to 1:
      $SplitTable[i] \leftarrow False$
      for j from i to n:
         If $IsWord(i,j)$ AND $SplitTable[j+1]$:
            $SplitTable[i] \leftarrow True$

   return $SplitTable[1]$

$n = 18$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k

$i = 14$

$j = 18$

$SplitTable$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k	
F	F	F	F	F	F	F	F	F	F	F	F	F	F	F	F	F	F	T

$i = 14$ $j + 1 = 19$

Dynamic Programming Approach

$FastSplittable(A[1..n])$:
    $SplitTable[n+1] \leftarrow True$

    for i from n to 1:
        $SplitTable[i] \leftarrow False$
        for j from i to n:
            If $IsWord(i,j)$ AND $SplitTable[j+1]$:
                $SplitTable[i] \leftarrow True$

    return $SplitTable[1]$

$n = 18$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k

$i = 14$

$j = 18$

$IsWord(14,18)$ is $True$!

$SplitTable$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k	
F	F	F	F	F	F	F	F	F	F	F	F	F	T	F	F	F	F	T

$i = 14$ $j + 1 = 19$

Dynamic Programming Approach

```
FastSplittable(A[1..n]):
   SplitTable[n + 1] ← True

   for i from n to 1:
      SplitTable[i] ← False
      for j from i to n:
         If IsWord(i, j) AND SplitTable[j + 1]:
            SplitTable[i] ← True

   return SplitTable[1]
```

$n = 18$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k

$i = 12$ $j = 13$

$IsWord(12,13)$ is $True$!

$SplitTable$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k	
F	F	F	F	F	F	F	F	F	F	F	T	F	T	F	F	F	F	T

$i = 12$ $j + 1 = 14$

Dynamic Programming Approach

```
FastSplittable(A[1..n]):
  SplitTable[n+1] ← True

  for i from n to 1:
    SplitTable[i] ← False
    for j from i to n:
      If IsWord(i,j) AND SplitTable[j+1]:
        SplitTable[i] ← True

  return SplitTable[1]
```

$n = 18$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k

$i = 10$ $j = 11$

$IsWord(10,11)$ is $True$!

$SplitTable$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k	
F	F	F	F	F	F	F	F	F	T	F	T	F	T	F	F	F	F	T

$i = 10$ $j + 1 = 12$

Dynamic Programming Approach

$n = 18$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k

$i = 9$ $j = 11$

```
FastSplittable(A[1..n]):
   SplitTable[n + 1] ← True

   for i from n to 1:
      SplitTable[i] ← False
      for j from i to n:
         If IsWord(i,j) AND SplitTable[j + 1]:
            SplitTable[i] ← True

   return SplitTable[1]
```

$IsWord(9,11)$ is $True$!

$SplitTable$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k	
F	F	F	F	F	F	F	F	T	T	F	T	F	T	F	F	F	F	T

$i = 9$ $j + 1 = 12$

Dynamic Programming Approach

$n = 18$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k

$i = 1$ $\uparrow$ $\uparrow$ $j = 3$

```
FastSplittable(A[1..n]):
   SplitTable[n + 1] ← True

   for i from n to 1:
      SplitTable[i] ← False
      for j from i to n:
         If IsWord(i, j) AND SplitTable[j + 1]:
            SplitTable[i] ← True

   return SplitTable[1]
```

$IsWord(1,3)$ is $True$!

$SplitTable$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k	
F	F	F	T	F	T	F	F	T	T	F	T	F	T	F	F	F	F	T

$i = 1$ $\uparrow$ $\uparrow$ $j + 1 = 4$

Dynamic Programming Approach

$n = 18$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k

$i = 1$ $j = 3$

```
FastSplittable(A[1..n]):
  SplitTable[n + 1] ← True

  for i from n to 1:
    SplitTable[i] ← False
    for j from i to n:
      If IsWord(i,j) AND SplitTable[j + 1]:
        SplitTable[i] ← True

  return SplitTable[1]
```

$IsWord(1,3)$ is $True$!

$SplitTable$

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k	
T	F	F	T	F	T	F	F	T	T	F	T	F	T	F	F	F	F	T

$i = 1$ $j + 1 = 4$

If T propagates all the way back to $i = 1$, we have a segmentation!

FastSplittable Analysis

t	h	e	b	r	o	w	n	f	o	x	i	s	q	u	i	c	k	
T	F	F	T	F	T	F	F	T	T	F	T	F	T	F	F	F	F	T

```
FastSplittable(A[1..n]):
   SplitTable[n + 1] ← True

   for i from n to 1:
      SplitTable[i] ← False
      for j from i to n:
         If IsWord(i,j) AND SplitTable[j + 1]:
            SplitTable[i] ← True

   return SplitTable[1]
```

If T propagates all the way back to $i = 1$, we have a segmentation!

Previously we had the recurrence:

$$T(n) = 2T(n-1) + c \leq O(2^n)$$

I argue we can just read off the running time of $FastSplittable$ from the pseudocode

Wrap up

Work on homework 2! Due Tuesday night at midnight.

Next time:

- Subset Sum revisited
- Edit Distance
- Knapsack problem

No new reading assignment (Chapter 3 Erickson)